

Almost all cancellative triple systems are tripartite

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Abstract

A triple system is cancellative if no three of its distinct edges satisfy $A \cup B = A \cup C$. It is tripartite if it has a vertex partition into three parts such that every edge has exactly one point in each part. It is easy to see that every tripartite triple system is cancellative. We prove that almost all cancellative triple systems with vertex set $[n]$ are tripartite. This sharpens a theorem of Nagle and Rödl [15] on the number of cancellative triple systems. It also extends recent work of Person and Schacht [16] who proved a similar result for triple systems without the Fano configuration.

Our proof uses the hypergraph regularity lemma of Frankl and Rödl [11], and a stability theorem for cancellative triple systems due to Keevash and the second author [12].

1 Introduction

Let F be a fixed graph or hypergraph. Say that a (hyper)graph is F -free if it contains no copy of F as a (not necessarily induced) sub(hyper)graph. Beginning with a result of Erdős-Kleitman-Rothschild [9], there has been much work concerning the number and structure of F -free graphs with vertex set $[n]$ (see, e.g. [8, 14, 17, 1, 2, 3, 5]). The strongest of these results essentially state that for a large class of graphs F , most of the F -free graphs with vertex set $[n]$ have a similar structure to the F -free graph with the maximum number of edges. Many of these results use the Szemerédi regularity lemma.

With the development of the hypergraph regularity Lemma, these problems can be attacked for hypergraphs. For brevity, we refer to a 3-uniform hypergraph as a triple system or 3-graph.

Definition. For a 3-graph F let $\text{Forb}(n, F)$ denote the set of (labeled) F -free 3-graphs on vertex set $[n]$.

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The first result in this direction was due to Nagle and Rödl [15] who proved that for a fixed 3-graph F ,

$$|Forb(n, F)| \leq 2^{\text{ex}(n, F) + o(n^3)},$$

where $\text{ex}(n, F)$ is the maximum number of edges in an F -free triple system on n vertices. Since there is no extremal result for hypergraphs similar to Turán's theorem for graphs, one cannot expect a general result that characterizes the structure of almost all F -free triple systems for large classes of F . Nevertheless, much is known about the extremal numbers for a few specific 3-graphs F and one could hope to obtain characterizations for these F . Recently, Person and Schacht [16] proved the first result of this kind, by showing that almost all triple systems on $[n]$ not containing a Fano configuration are 2-colorable. The key property that they used was the linearity of the Fano plane, namely the fact that every two edges of the Fano plane share at most one vertex. This enabled them to apply the (weak) 3-graph regularity lemma, which is almost identical to Szemerédi's regularity lemma. They then proved an embedding lemma for linear hypergraphs essentially following ideas from Kohayakawa-Nagle-Rödl-Schacht [13].

It is well-known that such an embedding lemma fails to hold for non-linear 3-graphs unless one uses the (strong) 3-graph regularity lemma, and operating in this environment is more complicated. In this paper, we address the situation for a particular non-linear F using this approach.

A triple system is tripartite or 3-partite if it has a vertex partition into three parts such that every edge has exactly one point in each part. Denote by $T(n)$ the number of 3-partite 3-graphs on $[n]$. Let

$$s(n) := \left\lfloor \frac{n}{3} \right\rfloor \left\lfloor \frac{n+1}{3} \right\rfloor \left\lfloor \frac{n+2}{3} \right\rfloor \sim \frac{n^3}{27}$$

be the maximum number of edges in a 3-partite triple system with n vertices. A triple system is cancellative if $A \cup B = A \cup C$ implies that $B = C$ for edges A, B, C . Every tripartite triple system is cancellative and Katona conjectured, and Bollobás [6] proved that the maximum number of edges in a cancellative triple system with n vertices is $s(n)$. It is easy to see that a cancellative triple system is one that contains no copy of

$$F_5 = \{123, 124, 345\} \quad \text{and} \quad K_4^- = \{123, 124, 234\}.$$

Later Frankl and Füredi [10] sharpened Bollobás' theorem by proving that $\text{ex}(n, F_5) = s(n)$ for $n > 3000$ (this was improved to $n > 33$ in [12]).

Our main result is the following.

Theorem 1. *Almost all F_5 -free 3-graphs on $[n]$ are 3-partite. More precisely there is a constant C such that*

$$|Forb(n, F_5)| < \left(1 + 2^{Cn - \frac{2n^2}{45}}\right) T(n). \quad (1)$$

Theorem 1 clearly implies the same result for cancellative 3-graphs which is stated in the abstract. As mentioned before, the proof of Theorem 1 uses the strong hypergraph regularity lemma, and stability theorems.

Using the fact that

$$\frac{4 \cdot 3^n}{n^2} 2^{s(n)} < T(n) < 3^n 2^{s(n)}$$

(see Lemma 5), we get the following improvement over the general result of Nagle and Rödl [15] which only implies that $|Forb(n, F_5)| < 2^{s(n)+o(n^3)}$.

Corollary 1. *As $n \rightarrow \infty$,*

$$\log_2 |Forb(n, F_5)| = s(n) + n \log_2 3 + \Theta(\log n).$$

In a forthcoming paper [4], we shall characterize the structure of almost all F -free 3-graphs, where $F = \{123, 124, 125, 345\}$. Note that such a fine statement as Theorem 1 is rare even for graphs: Prömel and Steger [17] characterized the structure of almost all F -free graphs when F has a color-critical edge, and Balogh, Bollobás and Simonovits [3] when $F = K(2, 2, 2)$.

2 Stability

The key idea in the proof of Theorem 1 is to reduce the problem to 3-graphs that are almost 3-partite. We associate a hypergraph with its edge set.

For a triple system \mathcal{H} with a 3-partition P of its vertices, say that an edge is *crossing* if it has exactly one point in each part, otherwise say that it is *non-crossing*. Let D_P be the set of non-crossing edges. An *optimal partition* $X \cup Y \cup Z$ of a triple system \mathcal{H} is a 3-partition of the vertices of \mathcal{H} which minimizes the number of non-crossing edges. Let $D = D_{\mathcal{H}}$ be the number of non-crossing (*bad*) edges in an optimal partition $X \cup Y \cup Z$. Define

$$Forb(n, F_5, \eta) := \{\mathcal{H} \subset [n]^3 : F_5 \not\subset \mathcal{H} \text{ and } D_{\mathcal{H}} \leq \eta n^3\}.$$

The first part of the proof of Theorem 1 is the following result, which we will prove in Section 4.

Theorem 2. *For every $\eta > 0$, there exists $\nu > 0$ and n_0 such that if $n > n_0$, then*

$$|Forb(n, F_5) - Forb(n, F_5, \eta)| < 2^{(1-\nu)\frac{n^3}{27}}.$$

3 Hypergraph Regularity

In this section, we quickly define the notions required to state the hypergraph regularity Lemma. These concepts will be used in Section 4 to prove Theorem 2. Further details can be found in [11] or [15]. As mentioned before we associate a hypergraph with its edge set.

A *k-partite cylinder* is a *k*-partite graph G with *k*-partition V_1, \dots, V_k , and we write $G = \cup_{i < j} G^{ij}$, where $G^{ij} = G[V_i \cup V_j]$ is the bipartite subgraph of G with parts V_i and V_j . For $B \in [k]^3$, the 3-partite cylinder $G(B) = \cup_{\{i,j\} \in [B]^2} G^{ij}$ is called a *triad*. For a 2-partite cylinder G , the *density* of the pair V_1, V_2 with respect to G is $d_G(V_1, V_2) = \frac{|G^{12}|}{|V_1||V_2|}$.

Given an integer $l > 0$ and real $\epsilon > 0$, a *k*-partite cylinder G is called an (l, ϵ, k) -*cylinder* if for every $i < j$, G^{ij} is ϵ -regular with density $1/l$. For a *k*-partite cylinder G , let $\mathcal{K}_3(G)$ denote the 3-graph on $V(G)$ whose edges correspond to triangles of G . An easy consequence of these definitions is the following fact.

Lemma 1. (Triangle Counting Lemma) *For integer $l > 0$ and real $\theta > 0$, there exists $\epsilon > 0$ such that every $(l, \epsilon, 3)$ -cylinder G with $|V_i| = m$ for all i satisfies*

$$|\mathcal{K}_3(G)| = (1 \pm \theta) \frac{m^3}{l^3}.$$

We now move on to 3-graph definitions. A *k-partite 3-cylinder* is a *k*-partite 3-graph \mathcal{H} with *k*-partition V_1, \dots, V_k . Here *k*-partite means that every edge of \mathcal{H} has at most one point in each V_i . Often we will say that these edges are crossing, and the edges that have at least two points in some V_i are non-crossing. Given $B \in [k]^3$, let $\mathcal{H}(B) = \mathcal{H}[\cup_{i \in B} V_i]$. Given *k*-partite cylinder G and *k*-partite 3-cylinder \mathcal{H} with the same vertex partition, say that G *underlies* \mathcal{H} if $\mathcal{H} \subset \mathcal{K}_3(G)$. In other words, \mathcal{H} consists only of triangles in G . Define the density $d_{\mathcal{H}}(G(B))$ of \mathcal{H} with respect to the triad $G(B)$ as the proportion of edges of \mathcal{H} on top of triangles of $G(B)$, if the latter quantity is positive, and zero otherwise. This definition leads to the more complicated definition of \mathcal{H} being (δ, r) -regular with respect to $G(B)$, where $r > 0$ is an integer and $\delta > 0$. If in addition $d_{\mathcal{H}}(G(B)) = \alpha \pm \delta$, then say that \mathcal{H} is (α, δ, r) -regular with respect to $G(B)$. We will not give the precise definitions of (α, δ, r) -regularity, and it suffices to take this definition as a “black box” that will be used later.

For a vertex set V , an (l, t, γ, ϵ) -partition \mathcal{P} of $[V]^2$ is a partition $V = V_0 \cup V_1 \cup \dots \cup V_t$ together with a collection of edge disjoint bipartite graphs P_a^{ij} , where $1 \leq i < j \leq t, 0 \leq a \leq l_{ij} \leq l$ that satisfy the following properties:

- (i) $|V_0| < t$ and $|V_i| = \lfloor \frac{n}{t} \rfloor := m$ for each $i > 0$,
- (ii) $\cup_{a=0}^{l_{ij}} P_a^{ij} = K(V_i, V_j)$ for all $1 \leq i < j \leq t$, where $K(V_i, V_j)$ is the complete bipartite graph with parts V_i, V_j ,
- (iii) all but $\gamma \binom{t}{2}$ pairs $\{v_i, v_j\}$, $v_i \in V_i, v_j \in V_j$, are edges of ϵ -regular bipartite graphs P_a^{ij} , and
- (iv) for all but $\gamma \binom{t}{2}$ pairs $\{i, j\} \in [t]^2$, we have $|P_0^{ij}| \leq \gamma m^2$ and $d_{P_a^{ij}}(V_i, V_j) = (1 \pm \epsilon) \frac{1}{l}$ for all $a \in [l_{ij}]$.

Finally, suppose that $\mathcal{H} \subset [n]^3$ is a 3-graph and \mathcal{P} is an (l, t, γ, ϵ) -partition of $[n]^2$ with

$m_{\mathcal{P}} = |V_1|$. For each triad $P \in \mathcal{P}$, let $\mu_P = \frac{|\mathcal{K}_3(P)|}{m_{\mathcal{P}}^3}$. Then \mathcal{P} is (δ, r) -regular if

$$\sum \{\mu_P : P \text{ is a } (\delta, r)\text{-irregular triad of } \mathcal{P}\} < \delta \left(\frac{n}{m_{\mathcal{P}}} \right)^3.$$

We can now state the Regularity Lemma due to Frankl and Rödl [11].

Theorem 3. (Regularity Lemma) *For every δ, γ with $0 < \gamma \leq 2\delta^4$, for all integers t_0, l_0 and for all integer-valued functions $r = r(t, l)$ and all functions $\epsilon(l)$, there exist T_0, L_0, N_0 such that every 3-graph $\mathcal{H} \subset [n]^3$ with $n \geq N_0$ admits a $(\delta, r(t, l))$ -regular $(l, t, \gamma, \epsilon(l))$ -partition for some t, l satisfying $t_0 \leq t < T_0$ and $l_0 \leq l < L_0$.*

To apply the Regularity Lemma above, we need to define a cluster hypergraph and state an accompanying embedding Lemma, sometimes called the Key Lemma. Given a 3-graph \mathcal{J} , let \mathcal{J}^2 be the set of pairs that lie in an edge of \mathcal{J} .

Cluster 3-graph. For given constants $k, \delta, l, r, \epsilon$ and sets $\{\alpha_B : B \in [k]^3\}$ of nonnegative reals, let \mathcal{H} be a k -partite 3-cylinder with parts V_1, \dots, V_k , each of size m . Let G be a graph, and $\mathcal{J} \subset [k]^3$ be a 3-graph such that the following conditions are satisfied.

- (i) $G = \cup_{\{i,j\} \in \mathcal{J}^2} G^{ij}$ is an underlying cylinder of \mathcal{H} such that for all $\{i, j\} \in \mathcal{J}^2$, G^{ij} is an $(l, \epsilon, 2)$ -cylinder.
- (ii) For each $B \in \mathcal{J}$, $\mathcal{H}(B)$ is (α_B, δ, r) -regular with respect to the triad $G(B)$.

Then we say that \mathcal{J} is the *cluster 3-graph* of \mathcal{H} .

Lemma 2. (Embedding Lemma) *Let $k \geq 4$ be fixed. For all $\alpha > 0$, there exists $\delta > 0$ such that for $l > \frac{1}{\delta}$, there exists r, ϵ such that the following holds: Suppose that \mathcal{J} is the cluster 3-graph of \mathcal{H} with underlying cylinder G and parameters $k, \delta, l, r, \epsilon, \{\alpha_B : B \in [k]^3\}$ where $\alpha_B \geq \alpha$ for all $B \in \mathcal{J}$. Then $\mathcal{J} \subset \mathcal{H}$.*

For a proof of the Embedding Lemma, see [15].

4 Most F_5 -free triple systems are almost tripartite

In this section we will prove Theorem 2. We will need the following stability result proved in [12]. The constants have been adjusted for later use.

Theorem 4. (Keevash-Mubayi [12]) *For every $\nu'' > 0$, there exist ν', t_2 such that every F_5 -free 3-graph on $t > t_2$ vertices and at least $(1 - 2\nu') \frac{t^3}{27}$ edges has a 3-partition for which the number of non-crossing edges is at most $\nu'' t^3$.*

Given $\eta > 0$, our constants will obey the following hierarchy:

$$\eta \gg \nu'' \gg \nu' \gg \nu \gg \sigma, \theta \gg \alpha_0, \frac{1}{t_0} \gg \delta \gg \gamma > \frac{1}{l_0} \gg \frac{1}{r}, \epsilon \gg \frac{1}{n_0}.$$

Before proceeding with further details regarding our constants, we define the *binary entropy function* $H(x) := -x \log_2 x - (1-x) \log_2 (1-x)$. We use the fact that for $0 < x < 0.5$ we have

$$\binom{n}{xn} < 2^{H(x)n}.$$

Additionally, if x is sufficiently small then

$$\sum_{i=0}^{xn} \binom{n}{i} < 2^{H(x)n}. \quad (2)$$

Detailed definition of constants.

Set

$$\nu'' = \frac{\eta}{1000} \quad (3)$$

and suppose that ν'_1 and t_2 are the outputs of Theorem 4 with input ν'' . Put

$$\nu' = \min\{\nu'_1, \nu''\} \quad \text{and} \quad \nu = (\nu')^4. \quad (4)$$

We choose

$$\theta = \frac{\nu}{4(1-\nu)}. \quad (5)$$

Choose σ_1 small enough so that

$$\left(1 - \frac{\nu}{2}\right) \frac{n^3}{27} + o(n^3) + H(\sigma)n^3 \leq \left(1 - \frac{\nu}{3}\right) \frac{n^3}{27} \quad (6)$$

holds for sufficiently large n . In fact the function denoted by $o(n^3)$ will actually be seen to be of order $O(n^2)$ so (6) will hold for sufficiently large n . Choose σ_2 small enough so that (2) holds for σ_2 . Let

$$\sigma = \min\{\sigma_1, \sigma_2\}.$$

Next we consider the Triangle Counting Lemma (Lemma 1) which provides an ϵ for each θ and l . Since θ is fixed, we may let $\epsilon_1 = \epsilon_1(l)$ be the output of Lemma 1 for each integer l .

For σ defined above, set

$$\delta_1 = \alpha_0 = \frac{\sigma}{100} \quad \text{and} \quad t_1 = \left\lceil \frac{1}{\delta_1} \right\rceil. \quad (7)$$

Let

$$t_0 = \max\{t_1, t_2, 33\}.$$

Now consider the Embedding Lemma (Lemma 2) with inputs $k = 5$ and α_0 defined above. The Embedding Lemma gives $\delta_2 = \delta_2(\alpha_0)$, and we set

$$\delta = \min\{\delta_1, \delta_2\}, \quad \gamma = \delta^4, \quad l_0 = \frac{2}{\delta}. \quad (8)$$

For each integer $l > \frac{1}{\delta}$, let $r = r(l)$ and $\epsilon_2 = \epsilon_2(l)$ be the outputs of Lemma 2. Set

$$\epsilon = \epsilon(l) = \min\{\epsilon_1(l), \epsilon_2(l)\}. \quad (9)$$

With these constants, the Regularity Lemma (Theorem 3) outputs N_0 . We choose n_0 such that $n_0 > N_0$ and every $n > n_0$ satisfies (6).

Proof of the Theorem 2.

We will prove that

$$|Forb(n, F_5) - Forb(n, F_5, \eta)| < 2^{(1-\frac{\nu}{3})\frac{n^3}{27}}.$$

This is of course equivalent to Theorem 2.

For each $\mathcal{G} \in Forb(n, F_5) - Forb(n, F_5, \eta)$, we use the Hypergraph Regularity Lemma, Theorem 3, to obtain a (δ, r) -regular (l, t, γ, ϵ) -partition $\mathcal{P} = \mathcal{P}_{\mathcal{G}}$. The input constants for Theorem 3 are as defined above and then Theorem 3 guarantees constants T_0, L_0, N_0 so that every 3-graph \mathcal{G} on $n > N_0$ vertices admits a (δ, r) -regular (l, t, γ, ϵ) -partition \mathcal{P} where $t_0 \leq t \leq T_0$ and $l_0 \leq l \leq L_0$. To this partition \mathcal{P} , associate a *density vector* $s = (s_{\{i,j,k\}_{a,b,c}})$ where $1 \leq i < j < k \leq t$ and $1 \leq a, b, c \leq l$ and

$$d_{\mathcal{G}}(P_a^{ij} \cup P_b^{jk} \cup P_c^{ik}) \in [s_{\{i,j,k\}_{a,b,c}} \delta, (s_{\{i,j,k\}_{a,b,c}} + 1) \delta].$$

For each $\mathcal{G} \in Forb(n, F_5, \eta)$, choose one (δ, r) -regular (l, t, γ, ϵ) -partition $\mathcal{P}_{\mathcal{G}}$ guaranteed by Theorem 3, and let $\mathcal{P} = \{\mathcal{P}_1, \dots, \mathcal{P}_p\}$ be the set of all such partitions over the family $Forb(n, F_5, \eta)$. Define an equivalence relation on $Forb(n, F_5, \eta)$ by letting $\mathcal{G} \sim \mathcal{G}'$ iff

- 1) $\mathcal{P}_{\mathcal{G}} = \mathcal{P}_{\mathcal{G}'}$ and
- 2) \mathcal{G} and \mathcal{G}' have the same density vector.

The number of equivalence classes q is the number of partitions times the number of density vectors. Consequently,

$$q \leq \left(\binom{T_0+1}{2} (L_0+1) \right)^{\binom{n}{2}} \left(\frac{1}{\delta} \right)^{(T_0+1)(L_0+1)^3} < 2^{O(n^2)}.$$

We will show that each equivalence class $C(\mathcal{P}_i, s)$ satisfies

$$|C(\mathcal{P}_i, s)| = 2^{(1-\frac{\nu}{2})\frac{n^3}{27} + H(\sigma)n^3}. \quad (10)$$

Combined with the upper bound for q and (6), we obtain

$$|Forb(n, F_5, \eta)| \leq 2^{O(n^2)} 2^{(1-\frac{\nu}{2})\frac{n^3}{27} + H(\sigma)n^3} \leq 2^{(1-\frac{\nu}{3})\frac{n^3}{27}}.$$

For the rest of the proof, we fix an equivalence class $C = C(\mathcal{P}, s)$ and we will show the upper bound in (10). We may assume that \mathcal{P} has vertex partition $[n] = V_0 \cup V_1 \cup \dots \cup V_t$, $|V_i| = m = \lfloor \frac{n}{t} \rfloor$ for all $i \geq 1$, and system of bipartite graphs P_a^{ij} , where $1 \leq i < j \leq t$, $0 \leq a \leq l_{ij} \leq l$.

Fix $\mathcal{G} \in C$. Let $\mathcal{E}_0 \subset \mathcal{G}$ be the set of triples that either

- (i) intersect V_0 , or
- (ii) have at least two points in some $V_i, i \geq 1$, or
- (iii) contain a pair in P_0^{ij} for some i, j , or
- (iv) contain a pair in some P_a^{ij} that is not ϵ -regular with density $\frac{1}{t}$.

Then

$$|\mathcal{E}_0| \leq tn^2 + t \left(\frac{n}{t}\right)^2 n + \gamma \binom{t}{2} n + 2\gamma \binom{t}{2} \left(\frac{n}{t}\right)^2 n.$$

Let $\mathcal{E}_1 \subset \mathcal{G} - \mathcal{E}_0$ be the set of triples $\{v_i, v_j, v_k\}$ such that either

- (i) the three bipartite graphs of \mathcal{P} associated with the pairs within the triple form a triad P that is not (δ, r) -regular with respect to $\mathcal{G}(\{i, j, k\})$, or
- (ii) the density $d_{\mathcal{G}}(P) < \alpha_0$.

Then

$$|\mathcal{E}_1| \leq 2\delta t^3 \left(\frac{n}{t}\right)^3 (1 + \theta) + \alpha_0 \binom{t}{3} l^3 \left(\frac{n}{t}\right)^3 \frac{1}{l^3}.$$

Let $\mathcal{E}_{\mathcal{G}} = \mathcal{E}_0 \cup \mathcal{E}_1$. Now (7) and (8) imply that

$$|\mathcal{E}_{\mathcal{G}}| \leq \sigma n^3.$$

Set $\mathcal{G}' = \mathcal{G} - \mathcal{E}_{\mathcal{G}}$.

Next we define $\mathcal{J}^C = \mathcal{J}^C(\mathcal{G}) \subset [t]^3 \times [l] \times [l] \times [l]$ as follows: For $1 \leq i < j < k \leq t$, $1 \leq a, b, c \leq l$, we have $\{i, j, k\}_{a,b,c} \in \mathcal{J}^C$ if and only if

- (i) $P = P_a^{ij} \cup P_b^{jk} \cup P_c^{ik}$ is an $(l, \epsilon, 3)$ -cylinder, and
- (ii) $\mathcal{G}'(\{i, j, k\})$ is $(\bar{\alpha}, \delta, r)$ -regular with respect to P , where $\bar{\alpha} \geq \alpha_0$.

We view \mathcal{J}^C as a multiset of triples on $[t]$. For each $\phi : \binom{[t]}{2} \rightarrow [l]$, let $\mathcal{J}_{\phi} \subset \mathcal{J}^C$ be the 3-graph on $[t]$ corresponding to the function ϕ (without parallel edges). In other words, $\{i, j, k\} \in \mathcal{J}_{\phi}$ iff the triples of \mathcal{G} that lie on top of the triangles of $P_a^{ij} \cup P_b^{jk} \cup P_c^{ik}$, $a = \phi(ij)$, $b = \phi(jk)$, $c = \phi(ik)$, are $(\bar{\alpha}, \delta, r)$ -regular and the underlying bipartite graphs $P_a^{ij}, P_b^{jk}, P_c^{ik}$ are all ϵ -regular with density $1/l$.

By our choice of the constants in (8) and (9), we see that any $\mathcal{F} \subset \mathcal{J}_{\phi}$ with five vertices is a cluster 3-graph for \mathcal{G} , and hence by the Embedding Lemma $\mathcal{F} \subset \mathcal{G}$. Since $F_5 \not\subset \mathcal{G}$, we conclude that $F_5 \not\subset \mathcal{J}_{\phi}$. It was shown in [12] that for $t \geq 33$, we have $\text{ex}(t, F_5) \leq \frac{t^3}{27}$. Since we know that $t \geq 33$, we conclude that

$$|\mathcal{J}_{\phi}| \leq \text{ex}(t, F_5) \leq \frac{t^3}{27}$$

for each $\phi : \binom{[t]}{2} \rightarrow [l]$. Recall from (4) that $\nu' = \nu^{1/4}$.

Lemma 3. Suppose that $|\mathcal{J}^C| > (1 - \nu) \frac{l^3 t^3}{27}$. Then for at least $(1 - \nu') l^{\binom{t}{2}}$ of the functions $\phi : \binom{[t]}{2} \rightarrow [l]$ we have

$$|\mathcal{J}_\phi| \geq (1 - \nu') \frac{|\mathcal{J}^C|}{l^3}.$$

Proof. Form the following bipartite graph: the vertex partition is $\Phi \cup \mathcal{J}^C$, where

$$\Phi = \left\{ \phi : \binom{[t]}{2} \rightarrow [l] \right\}$$

and the edges are of the form $\{\phi, \{i, j, k\}_{abc}\}$ if and only if $\phi \in \Phi$, $\{i, j, k\}_{abc} \in \mathcal{J}^C$ where $\phi(\{i, j\}) = a$, $\phi(\{j, k\}) = b$, $\phi(\{i, k\}) = c$. Let E denote the number of edges in this bipartite graph. Since each $\{i, j, k\}_{abc} \in \mathcal{J}^C$ has degree precisely $l^{\binom{t}{2}-3}$, we have

$$E = |\mathcal{J}^C| l^{\binom{t}{2}-3}.$$

Note that the degree of ϕ is $|\mathcal{J}_\phi|$. Suppose for contradiction that the number of ϕ for which $|\mathcal{J}_\phi| \geq (1 - \nu') \frac{|\mathcal{J}^C|}{l^3}$ is less than $(1 - \nu') l^{\binom{t}{2}}$. Then since $|\mathcal{J}_\xi| \leq \frac{t^3}{27}$ for each $\xi \in \Phi$, we obtain the upper bound

$$E \leq (1 - \nu') l^{\binom{t}{2}} \frac{t^3}{27} + \nu' l^{\binom{t}{2}} (1 - \nu') \frac{|\mathcal{J}^C|}{l^3}.$$

Dividing by $l^{\binom{t}{2}-3}$ then yields

$$|\mathcal{J}^C| \leq (1 - \nu') l^3 \frac{t^3}{27} + \nu' (1 - \nu') |\mathcal{J}^C|.$$

Simplifying, we obtain

$$(1 - \nu' (1 - \nu')) |\mathcal{J}^C| \leq (1 - \nu') l^3 \frac{t^3}{27}.$$

The lower bound $|\mathcal{J}^C| > (1 - \nu) \frac{l^3 t^3}{27}$ then gives

$$(1 - \nu' (1 - \nu')) (1 - \nu) < 1 - \nu'.$$

Since $\nu' = \nu^{1/4}$, the left hand side expands to

$$1 - \nu' + \nu^{1/2} - \nu + \nu^{5/4} - \nu^{3/2} > 1 - \nu'.$$

This contradiction completes the proof. \square

Claim 1.

$$|\mathcal{J}^C| \leq (1 - \nu) \frac{l^3 t^3}{27}.$$

Once we have proved Claim 1, the proof is complete by following the argument which is very similar to that in [15]. Define

$$S^C = \bigcup_{\{i,j,k\}_{abc} \in \mathcal{J}^C} \mathcal{K}_3(P_a^{ij} \cup P_b^{jk} \cup P_c^{ik}).$$

The Triangle Counting Lemma implies that $|\mathcal{K}_3(P_a^{ij} \cup P_b^{jk} \cup P_c^{ik})| < \frac{m^3}{l^3}(1 + \theta)$. Now Claim 1 and (5) give

$$|S^C| \leq \frac{m}{l^3}(1 + \theta)|\mathcal{J}^C| \leq m^3(1 + \theta)(1 - \nu)\frac{t^3}{27} < m^3\frac{t^3}{27}\left(1 - \frac{\nu}{2}\right) \leq \frac{n^3}{27}\left(1 - \frac{\nu}{2}\right).$$

Since $\mathcal{G}' \in S^C$ for every $\mathcal{G} \in C$,

$$|\{\mathcal{G}' : \mathcal{G} \in C\}| \leq 2^{(1 - \frac{\nu}{2})\frac{n^3}{27}}.$$

Each $\mathcal{G} \in C$ can be written as $\mathcal{G} = \mathcal{G}' \cup \mathcal{E}_{\mathcal{G}}$. In view of (2) and $|\mathcal{E}_{\mathcal{G}}| \leq \sigma n^3$, the number of $\mathcal{E}_{\mathcal{G}}$ with $\mathcal{G} \in C$ is at most $\sum_{i \leq \sigma n^3} \binom{n^3}{i} \leq 2^{H(\sigma)n^3}$. Consequently,

$$|C| \leq 2^{(1 - \frac{\nu}{2})\frac{n^3}{27} + H(\sigma)n^3}$$

and we are done.

Proof of Claim 1. Suppose to the contrary that $|\mathcal{J}^C| > (1 - \nu)\frac{l^3 t^3}{27}$. We apply Lemma 3 and conclude that for most functions ϕ the corresponding triple system \mathcal{J}_{ϕ} satisfies

$$|\mathcal{J}_{\phi}| \geq (1 - \nu')\frac{|\mathcal{J}^C|}{l^3} > (1 - \nu')(1 - \nu)\frac{t^3}{27} > (1 - 2\nu')\frac{t^3}{27}.$$

By Theorem 4, we conclude that for all of these ϕ , the triple system \mathcal{J}_{ϕ} has a 3-partition where the number of non-crossing edges is at most $\nu'' t^3$. We also conclude that the number of crossing triples that are not edges of \mathcal{J}_{ϕ} is at most

$$\left(\frac{2\nu'}{27} + \nu''\right)t^3 < \frac{5}{3}\nu''t^3. \quad (11)$$

Fix one such ϕ and let the optimal partition of \mathcal{J}_{ϕ} be $P_{\phi} = X \cup Y \cup Z$. Let $P = V_X \cup V_Y \cup V_Z$ be the corresponding vertex partition of $[n]$. In other words, V_X consists of the union of all those parts V_i for which $i \in X$ etc. We will show that P is a partition of $[n]$ where the number of non-crossing edges $|D_P|$ is fewer than ηn^3 . This contradicts the fact that $\mathcal{G} \in \text{Forb}(n, F_5) - \text{Forb}(n, F_5, \eta)$ and completes the proof of Theorem 2.

We have argued earlier that $|\mathcal{E}_{\mathcal{G}}| \leq \sigma n^3 \leq \frac{\eta}{2}n^3$ so it suffices to prove that $|D_P - \mathcal{E}_{\mathcal{G}}| \leq \frac{\eta}{2}n^3$.

Call a $\xi : \binom{[t]}{2} \rightarrow [l]$ *good* if it satisfies the conclusion of Lemma 3, otherwise call it *bad*. For each ξ and edge $\{i, j, k\} \in \mathcal{J}_{\xi}$, we have a, b, c defined by $a = \xi(\{i, j\})$ etc. let \mathcal{G}_{ξ} be the union, over all $\{i, j, k\} \in \mathcal{J}_{\xi}$, of the edges of \mathcal{G} that lie on top of the triangles in $P_a^{ij} \cup P_b^{jk} \cup P_c^{ik}$. Let D_{ξ} be the set of edges in \mathcal{G}_{ξ} that are non-crossing with respect to $P = V_X \cup V_Y \cup V_Z$. We will estimate $|D_P - \mathcal{E}_{\mathcal{G}}|$ by summing $|D_{\xi}|$ over all ξ . Please note that each $e \in D_P - \mathcal{E}_{\mathcal{G}}$ lies in exactly $l^{\binom{t}{2}-3}$ different D_{ξ} due to the definition of \mathcal{J}^C . Summing over all ξ gives

$$l^{\binom{t}{2}-3}|D_P - \mathcal{E}_{\mathcal{G}}| = \sum_{\xi: \binom{[t]}{2} \rightarrow [l]} |D_{\xi}| \leq \sum_{\xi \text{ good}} |D_{\xi}| + \sum_{\xi \text{ bad}} |D_{\xi}|.$$

Note that for a given edge $\{i, j, k\} \in \mathcal{J}_\phi$ the number of edges in \mathcal{G}_ϕ corresponding to this edge is the number of edges in $V_i \cup V_j \cup V_k$ on top of triangles formed by the three bipartite graphs, each of which is ϵ -regular of density $1/l$. By the Triangle Counting Lemma, the total number of such triangles is at most

$$2|V_i||V_j||V_k| \left(\frac{1}{l}\right)^3 < 2 \left(\frac{n}{t}\right)^3 \left(\frac{1}{l}\right)^3.$$

By Lemma 3, the number of bad ξ is at most $\nu' l^{\binom{t}{2}}$. So we have

$$\sum_{\xi \text{ bad}} |D_\xi| \leq \nu' l^{\binom{t}{2}} \binom{t}{3} 2 \left(\frac{n}{t}\right)^3 \left(\frac{1}{l}\right)^3 < \nu' l^{\binom{t}{2}-3} n^3.$$

It remains to estimate $\sum_{\xi \text{ good}} |D_\xi|$.

Fix a good ξ and let the optimal partition of \mathcal{J}_ξ be $P_\xi = A \cup B \cup C$ (recall that we know the number of non-crossing edges with respect to this partition is less than $\nu'' t^3$).

Claim 2. The number of crossing edges of P_ξ that are non-crossing edges of P_ϕ is at most $100\nu'' t^3$.

Suppose that Claim 2 was true. Then we would obtain

$$\sum_{\xi \text{ good}} |D_\xi| \leq l^{\binom{t}{2}} \left[100\nu'' t^3 \left(\frac{n}{t}\right)^3 \frac{2}{l^3} + \nu'' t^3 \left(\frac{n}{t}\right)^3 \frac{2}{l^3} \right] \leq l^{\binom{t}{2}-3} [202\nu'' n^3].$$

Explanation: We consider the contribution from the non-crossing edges of P_ϕ that are (i) crossing edges of P_ξ and (ii) non-crossing edges of P_ξ . We do not need to consider the contribution from the crossing edges of P_ϕ since by definition, these do not give rise to edges of D_P .

Altogether, using (3) we obtain

$$|D_P - \mathcal{E}_G| \leq (202\nu'' + \nu') n^3 < \frac{\eta}{2} n^3$$

and the proof is complete. We now prove Claim 2.

Proof of Claim 2. Suppose for contradiction that the number of crossing edges of P_ξ that are non-crossing edges of P_ϕ is more than $100\nu'' t^3$. Each of these edges intersects at most $3\binom{t}{2}$ other edges of \mathcal{J}_ξ , so by the greedy algorithm we can find a collection of at least $50\nu'' t$ of these edges that form a matching M . Pick one such edge $e = \{k, k', k''\} \in M$ and assume that k and k' lie in same part U of P_ϕ . Let d be the number of ways to choose a set of two triples $\{f, f'\}$ with $f = \{i, j, k\}$, $f' = \{i, j, k'\}$, $i, j \notin U \cup \{k''\}$ and i and j lie in distinct parts of P_ϕ . Since $|\mathcal{J}_\phi| > (1 - 2\nu') \frac{t^3}{27}$, $|D_{P_\phi}| \leq \nu'' t^3$ and ν', ν'' are sufficiently small

$$d \geq (\min\{|X|, |Y|, |Z|\} - 1)^2 \geq \frac{t^2}{10}.$$

As $\{e, f, f'\} \cong F_5$ there are at least d potential copies of F_5 that we can form using e and two crossing triples f, f' of P_ϕ . Suppose that $f = \{i, j, k\}, f' = \{i, j, k'\}$ are both in \mathcal{J}_ϕ for one such choice of $\{f, f'\}$. Consider the following eight bipartite graphs:

$$\begin{aligned} G^{ij} &= P_{\phi(\{i,j\})}^{ij}, & G^{jk} &= P_{\phi(\{j,k\})}^{jk}, & G^{ik} &= P_{\phi(\{i,k\})}^{ik}, & G^{jk'} &= P_{\phi(\{j,k'\})}^{jk'}, & G^{ik'} &= P_{\phi(\{i,k'\})}^{ik'}, \\ G^{kk'} &= P_{\xi(\{k,k'\})}^{kk'}, & G^{k'k''} &= P_{\xi(\{k',k''\})}^{k'k''}, & G^{kk''} &= P_{\xi(\{k,k''\})}^{kk''}. \end{aligned}$$

Set $G = \bigcup G^{uv}$ where the union is over the eight bipartite graphs defined above. Since $\{e, f, f'\} \subset \mathcal{J}_\phi \cup \mathcal{J}_\xi$, the 3-graph $J = \{e, f, f'\}$ associated with G and \mathcal{G} is a cluster 3-graph. By (8) and (9), we may apply the Embedding Lemma and obtain the contradiction $F_5 \subset \mathcal{G}$. We conclude that $f'' \notin \mathcal{J}_\phi$ for some $f'' \in \{f, f'\}$.

To each $e \in M$ we have associated at least d triples $f'' \notin \mathcal{J}_\phi$. Since M is a matching and $|e \cap f''| = 1$, each such f'' is counted at most three times. Summing over all $e \in M$, we obtain at least $\frac{|M|d}{3} \geq \frac{5}{3}\nu''t^3$ triples f'' that are crossing with respect to P_ϕ but are not edges of \mathcal{J}_ϕ . This contradicts (11) and completes the proof. \square

5 Proof of Theorem 1

In this section we complete the proof of Theorem 1. We begin with some preliminaries.

5.1 Inequalities

We shall use Chernoff's inequality as follows:

Theorem 5. *Let X_1, \dots, X_m be independent $\{0, 1\}$ random variables with $P(X_i = 1) = p$ for each i . Let $X = \sum_i X_i$. Then the following inequality holds for $a > 0$:*

$$P(X < \mathbb{E}X - a) < \exp(-a^2/(2pm)).$$

We will use the following easy statement.

Lemma 4. *Every graph G with n vertices contains a matching of size at least $\frac{|G|}{2n}$.*

Proof. Assume that there is a maximal matching of size r . The $2r$ vertices of the matching can cover at most $2rn$ edges, by the maximality of the matching there is no other edge in G . \square

Recall that $T(n)$ is the number of 3-partite 3-graphs with vertex set $[n]$ and $s(n) = \lfloor \frac{n+2}{3} \rfloor \cdot \lfloor \frac{n+1}{3} \rfloor \cdot \lfloor \frac{n}{3} \rfloor$. For a 3-partition A, B, C of a 3-graph, and $u \in A, v \in B$, write $L_C(u, v)$ or simply $L(u, v)$ for the set of $w \in C$ such that uvw is an edge. As usual, the multinomial coefficient $\binom{n}{a,b,c} = \frac{n!}{a!b!c!}$.

Lemma 5. *As $n \rightarrow \infty$ we have*

$$\left(\frac{1}{6} - o(1)\right) \binom{n}{\lfloor \frac{n+2}{3} \rfloor, \lfloor \frac{n+1}{3} \rfloor, \lfloor \frac{n}{3} \rfloor} 2^{s(n)} < T(n) < 3^n 2^{s(n)}. \quad (12)$$

In addition,

$$T(n-2) < \left(n^2 2^{-\frac{2n^2}{9} + n}\right) T(n). \quad (13)$$

Proof. For the upper bound in (12), observe that 3^n counts the number of 3-partitions of the vertices, and the exponent is the maximum number of crossing edges that a 3-partite 3-graph can have.

For the lower bound we count the number of (unordered) 3-partitions where this equality can be achieved. Each such 3-partition gives rise to $2^{s(n)}$ 3-partite 3-graphs. The number of such 3-partitions of $[n]$ is at least

$$\frac{1}{6} \binom{n}{\lfloor \frac{n+2}{3} \rfloor, \lfloor \frac{n+1}{3} \rfloor, \lfloor \frac{n}{3} \rfloor}.$$

We argue next that most of the 3-partite 3-graphs obtained in this way are different. More precisely, we show below that for any given 3-partition P as above, most 3-partite 3-graphs with 3-partition P have a unique 3-partition (which must be P). Given a 3-partition U_1, U_2, U_3 of $[n]$, if the crossing edges are added randomly, then Chernoff's inequality gives that almost all 3-graphs generated satisfy the following two conditions:

- (i) for all $u \in U_i, v \in U_j$, where $\{i, j, \ell\} = \{1, 2, 3\}$ we have $|L_{U_\ell}(u, v)| > n/10$
- (ii) for $\{i, j, \ell\} = \{1, 2, 3\}$ and for every $A_i \subset U_i, A_j \subset U_j$ with $|A_i|, |A_j| > n/10$ and $v \in U_\ell$, the number of crossing edges intersecting each of A_i, A_j and containing v is at least $|A_i||A_j|/10$.

If \mathcal{H} has 3-partition U_1, U_2, U_3 of $[n]$, and it satisfies conditions (i) and (ii), then the 3-partition is unique. Indeed, take u, v lying in an edge, then u, v and $L(u, v)$ are in different parts, where $|L(u, v)| > n/10$, so for $w \in L(u, v)$, $L(u, w)$ is in the same part as v and $L(v, w)$ is in the same part as u . Now by (ii) the rest of the vertices must lie in a unique part.

To prove (13) first note that if $a + b + c = n$, then $\binom{n}{a, b, c}$ is maximized for $a = \lfloor (n+2)/3 \rfloor, b = \lfloor (n+1)/3 \rfloor, c = \lfloor n/3 \rfloor$. This implies that

$$3^n = \sum_{a+b+c=n} \binom{n}{a, b, c} \leq \binom{n+2}{2} \binom{n}{\lfloor \frac{n+2}{3} \rfloor, \lfloor \frac{n+1}{3} \rfloor, \lfloor \frac{n}{3} \rfloor} < (0.6)n^2 \binom{n}{\lfloor \frac{n+2}{3} \rfloor, \lfloor \frac{n+1}{3} \rfloor, \lfloor \frac{n}{3} \rfloor}.$$

Together with (12) we obtain

$$\frac{T(n-2)}{T(n)} < \frac{3^{n-2} 2^{s(n-2)}}{\left(\frac{1}{6} - o(1)\right) \binom{n}{\lfloor \frac{n+2}{3} \rfloor, \lfloor \frac{n+1}{3} \rfloor, \lfloor \frac{n}{3} \rfloor} 2^{s(n)}} < n^2 2^{s(n-2) - s(n)}.$$

It is easy to see that $s(n) - s(n-2) \geq 2n^2/9 - n$, and the result follows. \square

5.2 Lower Density

Definition 6. A vertex partition U_1, U_2, U_3 of a 3-graph \mathcal{F} is μ -lower dense if each of the following conditions are satisfied:

(i) For every i if $A_i \subset U_i$ with $|A_i| \geq \mu n$ then

$$|\{E \in \mathcal{F} : |E \cap A_i| = 1, \text{ for } 1 \leq i \leq 3\}| > |A_1| \cdot |A_2| \cdot |A_3| \cdot 2^{-3}.$$

(ii) Let $\{i, j, \ell\} = \{1, 2, 3\}$, $A_i \subset U_i$ with $|A_i| \geq \mu n$, $G \subset U_j \times U_\ell$ with $|G| \geq \mu^2 n^2$. Then

$$|\{E \in \mathcal{F} : |E \cap A_i| = 1, E - A_i \in G\}| > |A_i| \cdot |G| \cdot 2^{-3}.$$

(iii) Let $\{i, j, \ell\} = \{1, 2, 3\}$, $A_i \subset U_i$ and $A_j \subset U_j$ with $|A_i|, |A_j| \geq \mu n$, and G be a matching on U_ℓ with $|G| \geq \mu n$. Set

$$\mathcal{F}_{A_i, A_j, G} = \{\{C, D\} \in \mathcal{F}^2 : C - U_\ell = D - U_\ell, |C \cap A_i| = |C \cap A_j| = 1, \{(C \cap U_\ell), (D \cap U_\ell)\} \in G\}.$$

Then

$$|\mathcal{F}_{A_i, A_j, G}| \geq \frac{|A_i| \cdot |A_j| \cdot |G|}{2^7}.$$

(iv) For every i we have $||U_i| - n/3| < \mu n$.

For $\mu > 0$ let $\text{Forb}(n, F_5, \eta, \mu) \subset \text{Forb}(n, F_5, \eta)$ be the family of μ -lower dense hypergraphs.

Lemma 6. For every η if $\mu^3 \geq 10^3 H(6\eta)$ then for n large enough

$$|\text{Forb}(n, F_5, \eta) - \text{Forb}(n, F_5, \eta, \mu)| < 2^{n^3(1/27 - \mu^3/40)}.$$

Proof. We wish to count the number of $\mathcal{H} \in \text{Forb}(n, F_5, \eta) - \text{Forb}(n, F_5, \eta, \mu)$. The number of ways to choose a 3-partition of \mathcal{H} is at most 3^n . Given a particular 3-partition $P = (U_1, U_2, U_3)$, the number of ways the at most ηn^3 bad edges could be placed is at most

$$\sum_{i \leq \eta n^3} \binom{\binom{n}{3}}{i} < 2^{H(6\eta) \binom{n}{3}}.$$

If $|U_i - n/3| > \mu n$ for some i , then the number of possible crossing edges is at most

$$n^3(1/27 - \mu^2/4 + \mu^3/4) < n^3(1/27 - \mu^2/5).$$

We conclude that the number of $\mathcal{H} \in \text{Forb}(n, F_5, \eta) - \text{Forb}(n, F_5, \eta, \mu)$ for which there exists a partition that fails property (iv) is at most

$$f(n, \eta) 2^{n^3(1/27 - \mu^2/5)},$$

where

$$f(n, \eta) = 3^n \cdot 2^{H(6\eta) \binom{n}{3}}.$$

Since $\mathcal{H} \notin \text{Forb}(n, F_5, \eta, \mu)$ it fails to satisfy one of the four conditions in Definition 6. For a fixed partition P and choice of bad edges, we may view \mathcal{H} as a probability space where we choose each crossing edge with respect to P independently with probability $1/2$. The total number of ways to choose the crossing edges is at most $2^{n^3/27}$ (an upper bound on the size of the probability space) so we obtain that $|\text{Forb}(n, F_5, \eta) - \text{Forb}(n, F_5, \eta, \mu)|$ is upper bounded by

$$f(n, \eta) \cdot 2^{n^3/27} \cdot \text{Prob}(\mathcal{H} \text{ fails (i) or (ii) or (iii)}) + f(n, \eta) 2^{n^3(1/27 - \mu^2/5)}.$$

We will consider each of these probabilities separately and then use the union bound. First however, note that the number of choices for $A_i \subset U_i$ as in Definition 6 is at most 2^n and the number of ways G could be chosen is at most 2^{n^2} .

(i) Since $|A_1||A_2||A_3| \geq \mu^3 n^3$, Chernoff's inequality gives

$$\text{Prob}(\mathcal{H} \text{ fails (i)}) \leq 2^{3n} \cdot \exp(-\mu^3 n^3/16).$$

(ii) Since $|A_i||G| \geq \mu^3 n^3$, Chernoff's inequality gives

$$\text{Prob}(\mathcal{H} \text{ fails (ii)}) \leq 2^n \cdot 2^{n^2} \cdot \exp(-\mu^3 n^3/16).$$

(iii) Since $|A_i||A_j||G| \geq \mu^3 n^3$ and both edges C and D must be present, we apply Chernoff's inequality with $m = |A_i||A_j||G|/2$ and $p = 1/4$. The number of matchings G is at most $(n^2)^{n/2} = 2^{n \log_2 n}$, so

$$\text{Prob}(\mathcal{H} \text{ fails (iii)}) \leq 2^{2n} \cdot 2^{n \log_2 n} \cdot \exp(-\mu^3 n^3/32).$$

The lemma now follows since $10^3 H(6\eta) \leq \mu^3$, and n is sufficiently large. \square

5.3 There is no bad vertex

Let $\mathcal{H} \in \text{Forb}(n, F_5, \eta, \mu)$, assume n is large enough, and U_1, U_2, U_3 is an optimal partition of \mathcal{H} , with $x \in U_1$. For a vertex y let $L_{i,j}(y)$ denote the set of edges of \mathcal{H} containing y , and additionally intersecting U_i and U_j . In particular, $L_{i,i}(y)$ is the set of edges of \mathcal{H} which contain y , and their other vertices are in U_i .

The aim of this subsection is to prove the following lemma, which shows that the number of bad edges containing a vertex is small.

Lemma 7. *Each of the followings is satisfied for $x \in U_1$.*

- (i) $|L_{1,1}(x)| < 2\mu n^2$.
- (ii) $|L_{1,2}(x)| < 2\mu n^2$.
- (iii) $|L_{2,2}(x)| < 2\mu n^2$.
- (iv) $|L_{1,3}(x)| < 2\mu n^2$.
- (v) $|L_{3,3}(x)| < 2\mu n^2$.

Proof. (i) If $|L_{1,1}(x)| > 2\mu n^2$ then by Lemma 4 $\{E - x : E \in L_{1,1}(x)\}$ contains a matching G with size at least μn . Then using Definition 6 (iii) (with $G, A_i = U_2, A_j = U_3$) for \mathcal{H} , we find $y, z \in U_1, a \in U_2, b \in U_3$ such that $xyz, yab, zab \in \mathcal{H}$, yielding an $F_5 \subset \mathcal{H}$, a contradiction.

(ii) Suppose for contradiction that $|L_{1,2}(x)| \geq 2\mu n^2$. By the optimality of the partition $|L_{1,2}(x)| \leq |L_{2,3}(x)|$, otherwise x could be moved to U_3 to decrease the number of bad edges. We shall use property (ii) in Definition 6. We use it with $G = \{E - x : E \in L_{1,2}(x)\}$ and

$$A_3 = \{z \in U_3 : \exists \text{ crossing edges } E_1, E_2 \in \mathcal{H} \text{ with } \{x, z\} \subset E_1 \cap E_2\}.$$

Note that $|A_3| \geq \mu n$ as $|L_{2,3}(x)| \geq 2\mu n^2$. Since \mathcal{H} is μ -lower dense, we find $abz \in \mathcal{H}$ with $xab \in L_{1,2}(x)$ and $z \in A_3$. By definition of A_3 , there exists $b' \in U_2 - \{b\}$ such that $xb'z \in L_{2,3}(x)$. This gives us $abx, abz, xb'z \in \mathcal{H}$, forming an F_5 .

(iii) Suppose for contradiction that $|L_{2,2}(x)| \geq 2\mu n^2$. By Lemma 4 $\{E - x : E \in L_{2,2}(x)\}$ contains a matching G with size at least μn . Then using Definition 6 (iii) (with $G, A_i = U_1 - x, A_j = U_3$) we find $b, b' \in U_2, a \in U_1, c \in U_3$ such that $abc, ab'c, xbb' \in \mathcal{H}$, forming an $F_5 \subset \mathcal{H}$, a contradiction.

The proof of (iv) is identical to (ii) and of (v) is to (iii). \square

5.4 Getting rid of bad edges - A Progressive Induction

Here we have to do something similar to the previous section, however, as we get rid of only a few edges, the computation needed is more delicate. We shall do progressive induction on the number of vertices. The general idea is that we remove some vertices of a bad edge, and count the number of ways it could have been joined to the rest of the hypergraph.

We shall prove (1) via induction on n . Fix an n_0 such that $1/n_0$ is much smaller than any of our constants, and all of our prior lemmas and theorems are valid for every $n \geq n_0$. Let $C > 10$ be sufficiently large that (1) is true for every $n \leq n_0$.

Let $Forb'(n, F_5, \eta, \mu)$ be the set of hypergraphs $\mathcal{H} \in Forb(n, F_5, \eta, \mu)$ having an optimal partition with a bad edge. Our final step is to give an upper bound $|Forb'(n, F_5, \eta, \mu)|$. There are two types of bad edges, one which is completely inside of a class, and the one which intersects two classes.

Let the bad edge be xyz , and the optimal partition be U_1, U_2, U_3 . Without loss of generality assume that $x, y \in U_1$.

In an $\mathcal{H} \in Forb'(n, F_5, \eta, \mu)$, x, y, z could be chosen at most n^3 ways, the optimal partition of \mathcal{H} in at most 3^n ways and the hypergraph $\mathcal{H} - \{x, y\}$ in at most $|Forb(n - 2, F_5)|$ ways. By Lemma 7 each of $|L_{1,1}(x)|, |L_{1,1}(y)|, |L_{1,2}(x)|, |L_{1,2}(y)|, |L_{1,3}(x)|, |L_{1,3}(y)|, |L_{2,2}(x)|, |L_{2,2}(y)|, |L_{3,3}(x)|, |L_{3,3}(y)|$ is at most $2\mu n^2$, therefore the number of ways the bad edges

could be joined to x, y is at most

$$\left(\sum_{i \leq 2\mu n^2} \binom{n^2/2}{i} \right)^{10} \leq 2^{10H(4\mu)n^2}.$$

The key point is that for any $(u, v) \in (U_2 - z) \times (U_3 - z)$, we cannot have both $xuv, yuv \in \mathcal{H}$ otherwise they form with xyz a copy of F_5 . Together with Definition 6 part (iv), we conclude that the number of ways to choose the crossing edges containing x or y is at most

$$3^{|U_2||U_3|} 2^{2n} \leq 3^{\frac{n^2}{9} + \mu n^2}.$$

Note that the 2^{2n} estimates the number of ways having edges containing u, z or vz , as for these pairs we do not have any restriction.

Putting this together,

$$|Forb'(n, F_5, \eta, \mu)| \leq n^3 3^n |Forb(n-2, F_5)| \cdot 2^{10H(4\mu)n^2} 3^{\frac{n^2}{9} + \mu n^2}. \quad (14)$$

By the induction hypothesis, this is at most

$$n^3 3^n (1 + 2^{C(n-2) - \frac{2(n-2)^2}{45}}) T(n-2) 2^{10H(4\mu)n^2} 3^{\frac{n^2}{9} + \mu n^2}.$$

Using (13) this is upper bounded by

$$n^5 3^n \left(1 + 2^{C(n-2) - \frac{2(n-2)^2}{45}} \right) 2^{(90H(4\mu) + \log_2 3 + 9\mu - 2 + \frac{9}{n}) \frac{n^2}{9}} \cdot T(n).$$

As mentioned before, the crucial point in the expression above is that $\log_2 3 - 2 < 0$. More precisely, since $n > n_0$, $\log_2 3 < 1.59$ and $90H(4\mu) + 9\mu < 0.001$, we have

$$\left(90H(4\mu) + \log_2 3 + 9\mu - 2 + \frac{9}{n} \right) \frac{n^2}{9} < -\frac{2n^2}{45}.$$

Consequently,

$$|Forb'(n, F_5, \eta, \mu)| \leq n^5 3^n \left(1 + 2^{C(n-2) - \frac{2(n-2)^2}{45}} \right) \cdot 2^{-\frac{2n^2}{45}} T(n) < \frac{1}{10} 2^{Cn - \frac{2n^2}{45}} T(n).$$

Now we can complete the proof of (1) by upper bounding $|Forb(n, F_5)|$ as follows:

$$\begin{aligned} & |Forb(n, F_5) - Forb(n, F_5, \eta)| + |Forb(n, F_5, \eta) - Forb(n, F_5, \eta, \mu)| + |Forb'(n, F_5, \eta, \mu)| + T(n) \\ & < 2^{(1-\nu)\frac{n^3}{27}} + 2^{n^3(\frac{1}{27} - \frac{\mu^3}{40})} + \frac{1}{10} 2^{Cn - \frac{2n^2}{45}} T(n) + T(n) \\ & < (1 + 2^{Cn - \frac{2n^2}{45}}) T(n), \end{aligned}$$

where the last inequality holds due to $T(n) > 2^{s(n)} > 2^{\frac{n^3}{27} - O(n^2)}$. This completes the proof of the theorem. \square

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